



INSTITUT
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AN EVENT-B MODEL OF A MECHANICAL LUNG VENTILATOR

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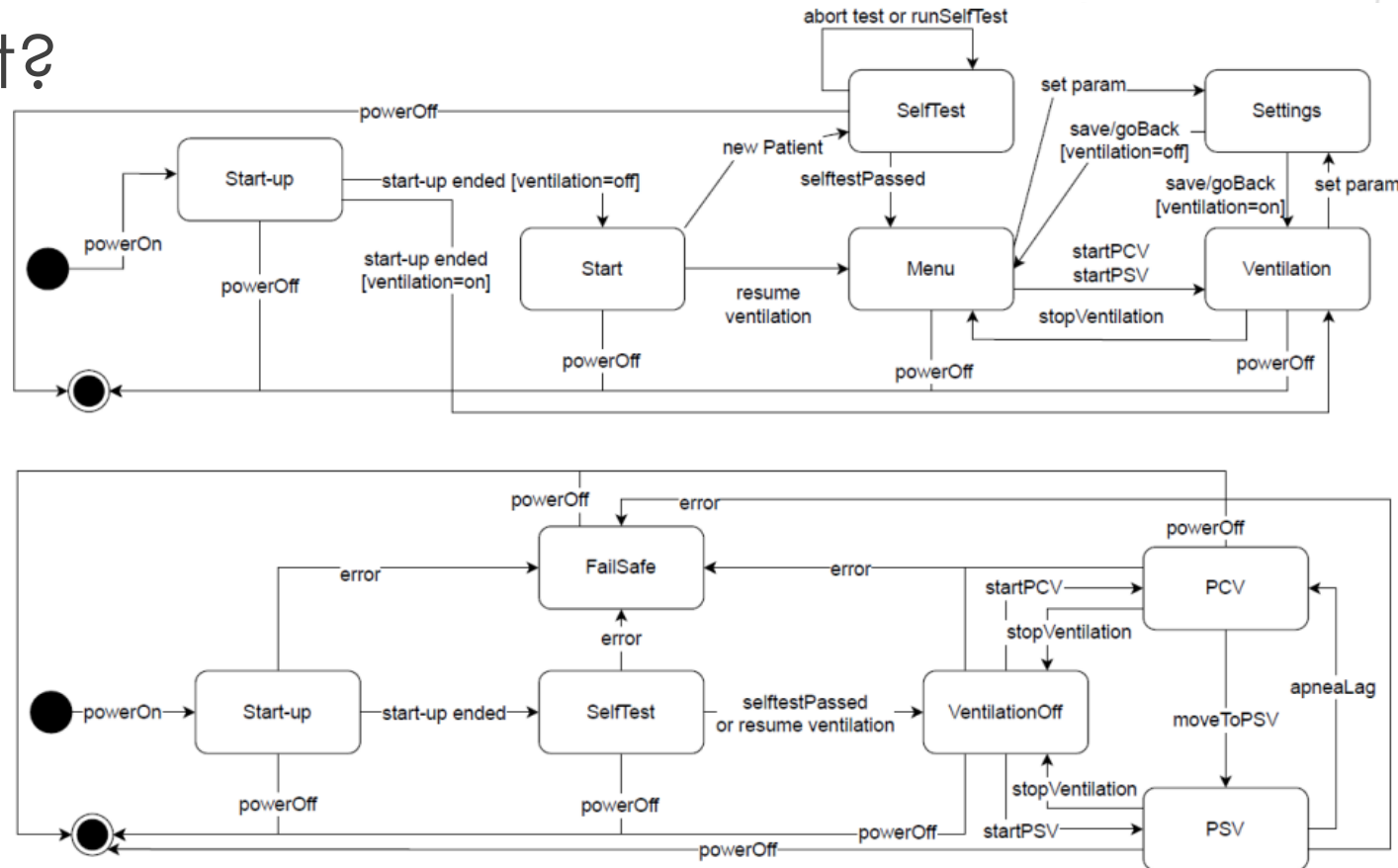


EVENT-B METHOD

- ❑ State-based formal method
- ❑ A model is made up on:
 - Contexts: sets (user types), constants and axioms
 - Machines/Refinements: variables, invariants, events
- ❑ Proof correctness:
 - Machine/Refinement: each event must preserve the invariant of the machine/refinement
 - Refinement:
 - The guard of a concrete event must be stronger than that of the abstract one
 - The effect of a concrete event is included in that of the abstract one

MODELING STRATEGY: 1/4

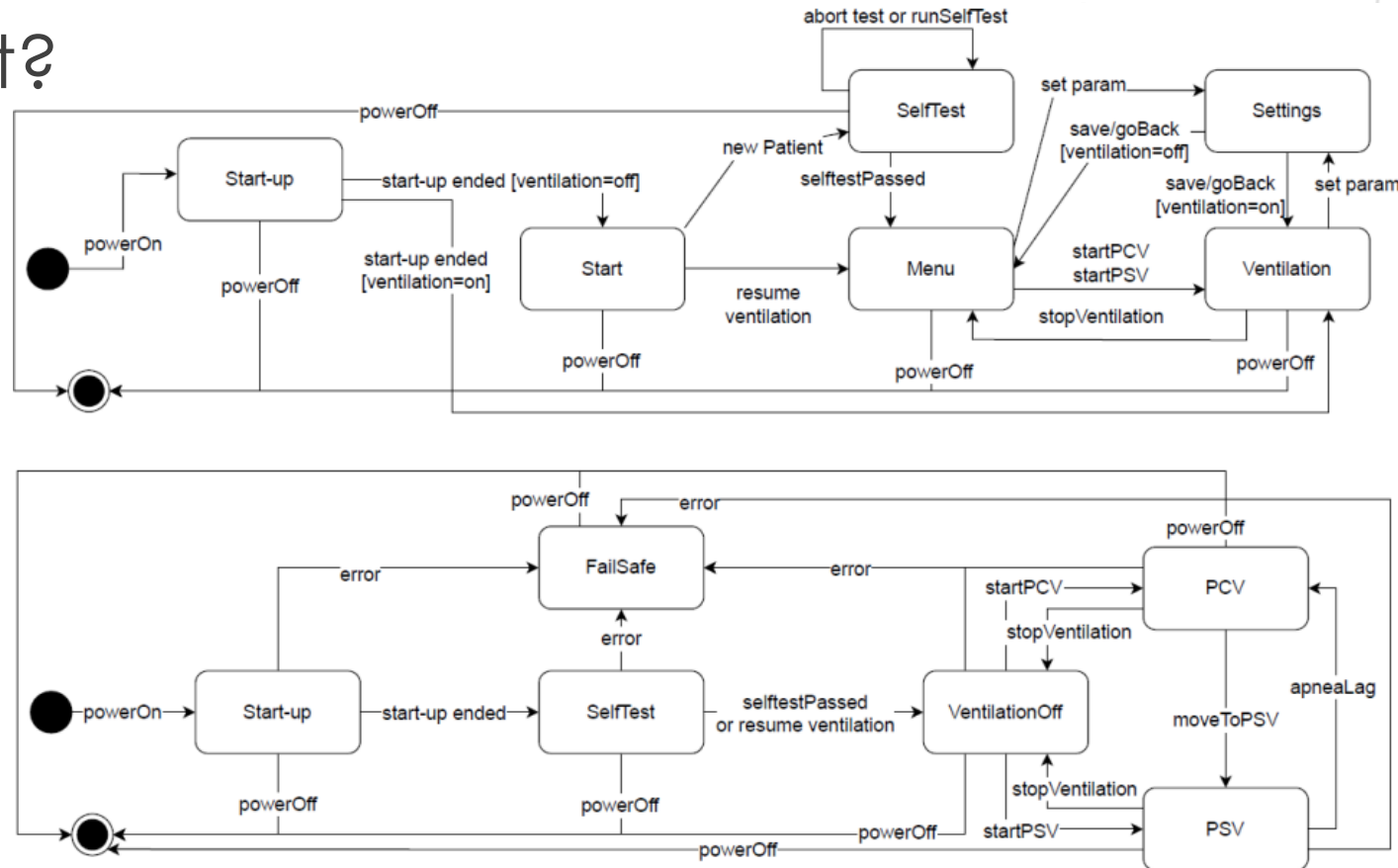
Which component, GUI or the controller, should be modelled at first?



First step of the controller's behavior is independent from that of the GUI

MODELING STRATEGY: 1/4

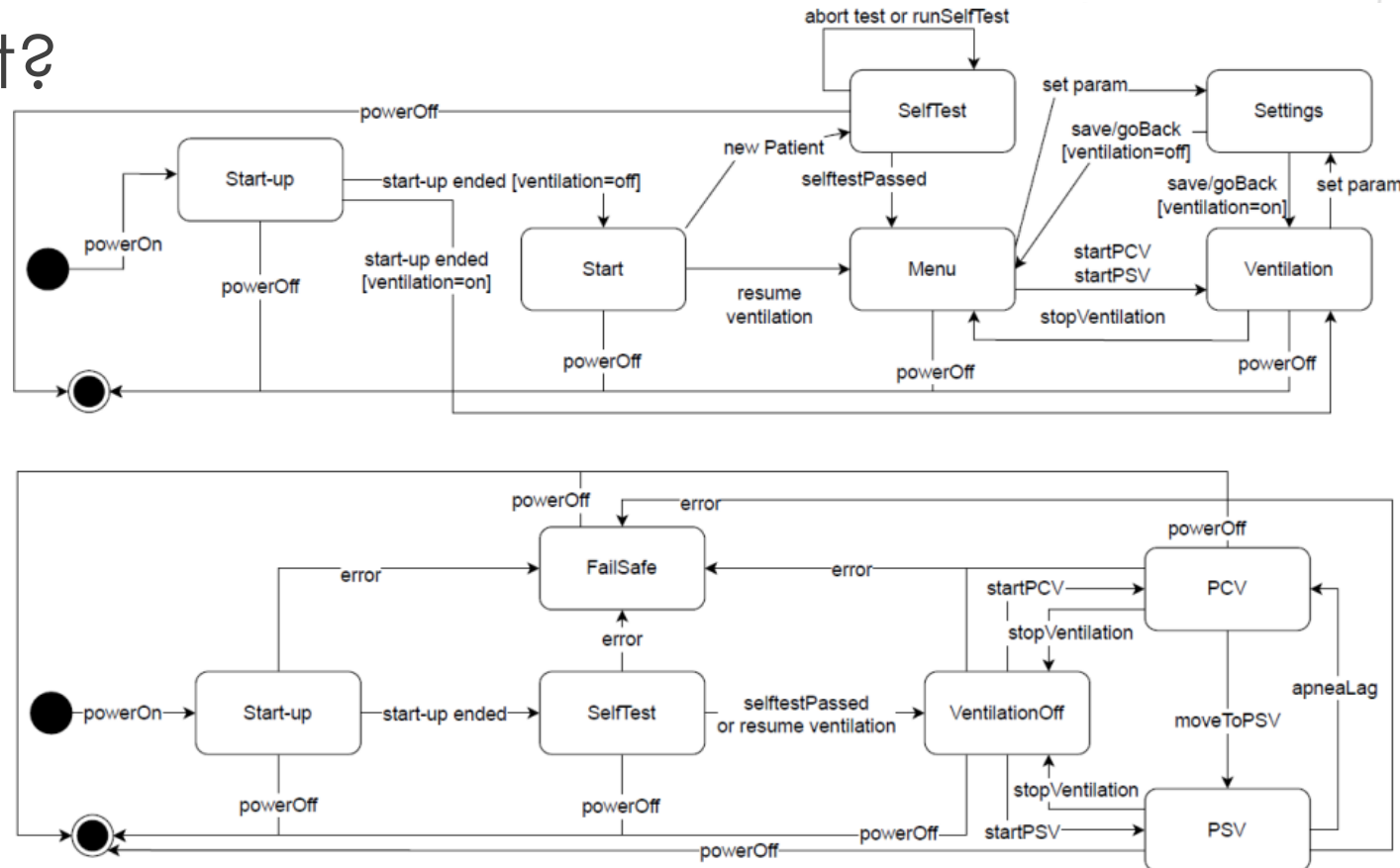
□ Which component, GUI or the controller, should be modelled at first?



Main functionalities of the controller depend on the orders received from the GUI

MODELING STRATEGY: 1/4

Which component, GUI or the controller, should be modelled at first?



GUI functionalities is modeled before those of the controller

MODELING STRATEGY: 2/4

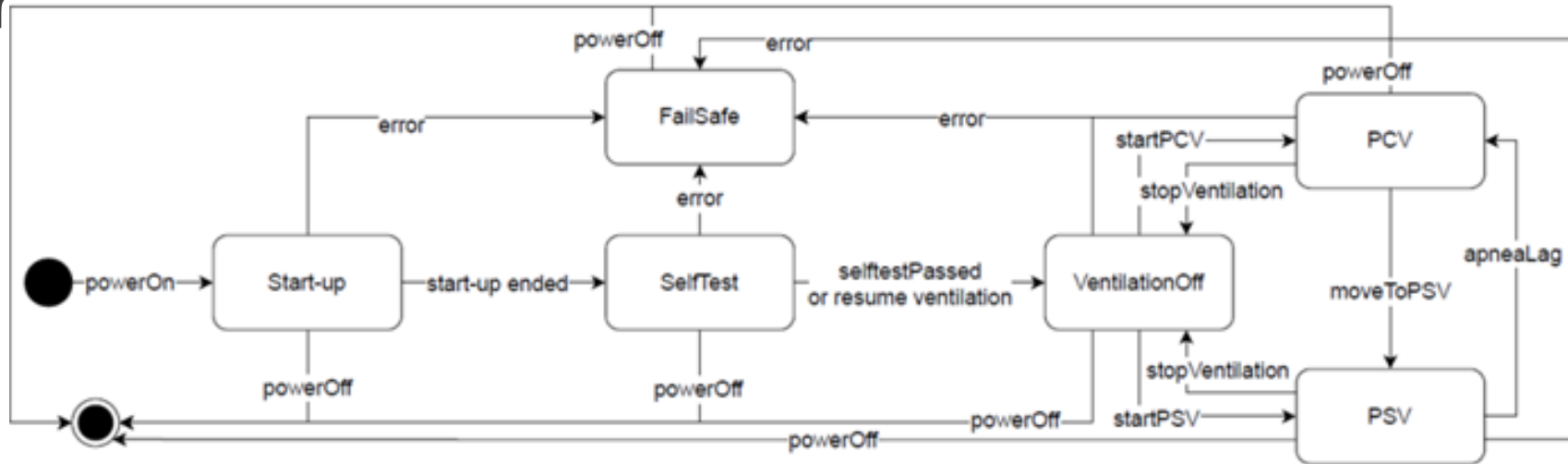
□ Which level is the best to introduce time?

The behavior of the system may depend on the internal battery whose level depends on time progression

The level of the battery is modelled in an early level with the usual time progress event

MODELING STRATEGY: 3/4

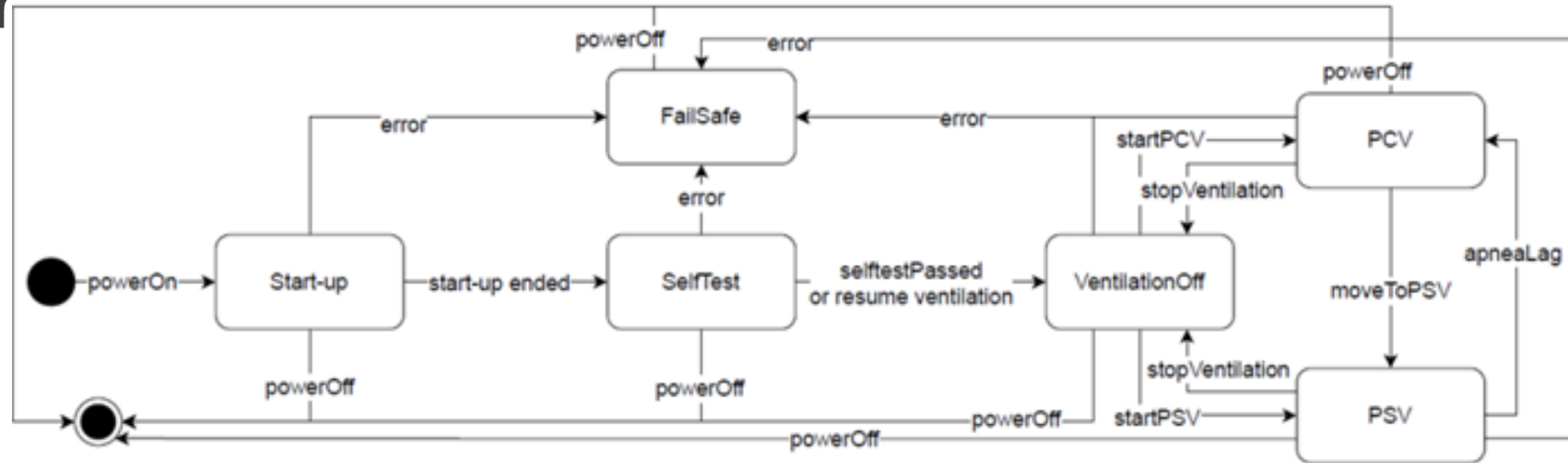
- Can the controller event error be dealt with like other ever



- an error may be due to several causes (valve failure, backup battery failure, etc.)
- an error may be raise at any moment

MODELING STRATEGY: 3/4

- Can the controller event error be dealt with like other ever



Event error is modeled as a refinement of the event progress

MODELING STRATEGY: 4/4

- ❑ Can the transitions common to the GUI and the controller (powerOn, powerOff, start-up ended, etc.) be represented by a single event?

Transitions with the same effect: a single Event-B event (powerOn, powerOff)

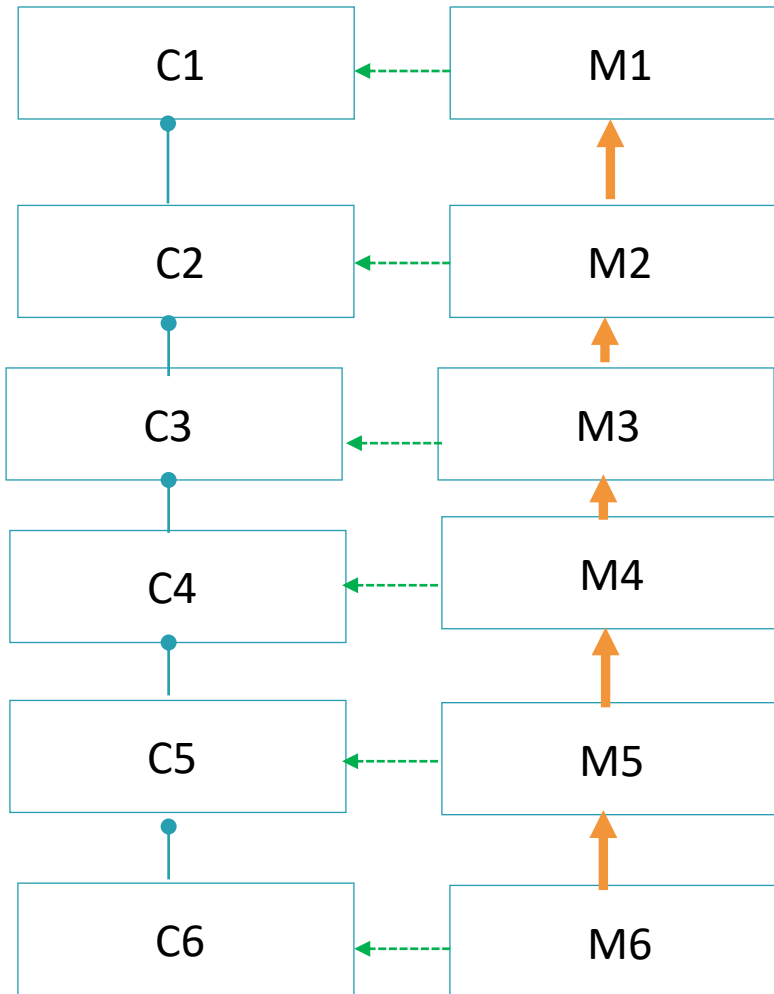
Transitions with different effects: distinct Event-B events (start-up ended)

EVENT-B MODEL STRUCTURE

—●— EXTENDS

- - - -> SEES

—> REFINES



GUI functionalities

Controller functionalities

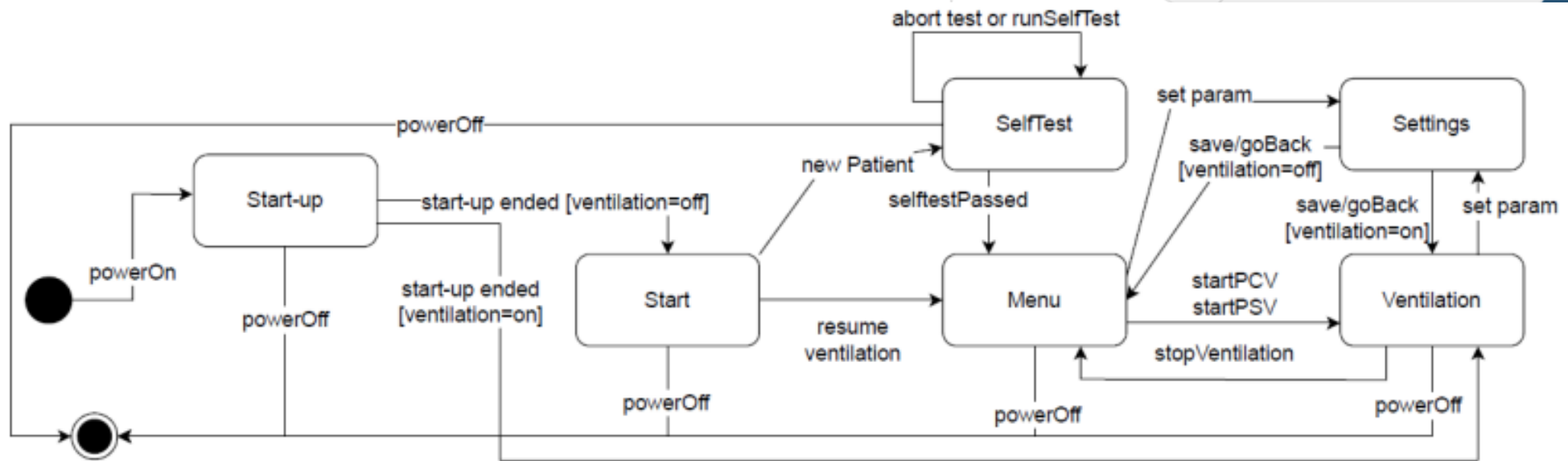
Ventilation modes: PCV, PSV

Ventilation phases: inspiration, expiration, etc.

Valves

Alarms

1ST LEVEL: GUI FUNCTIONALITIES (M1+C1)

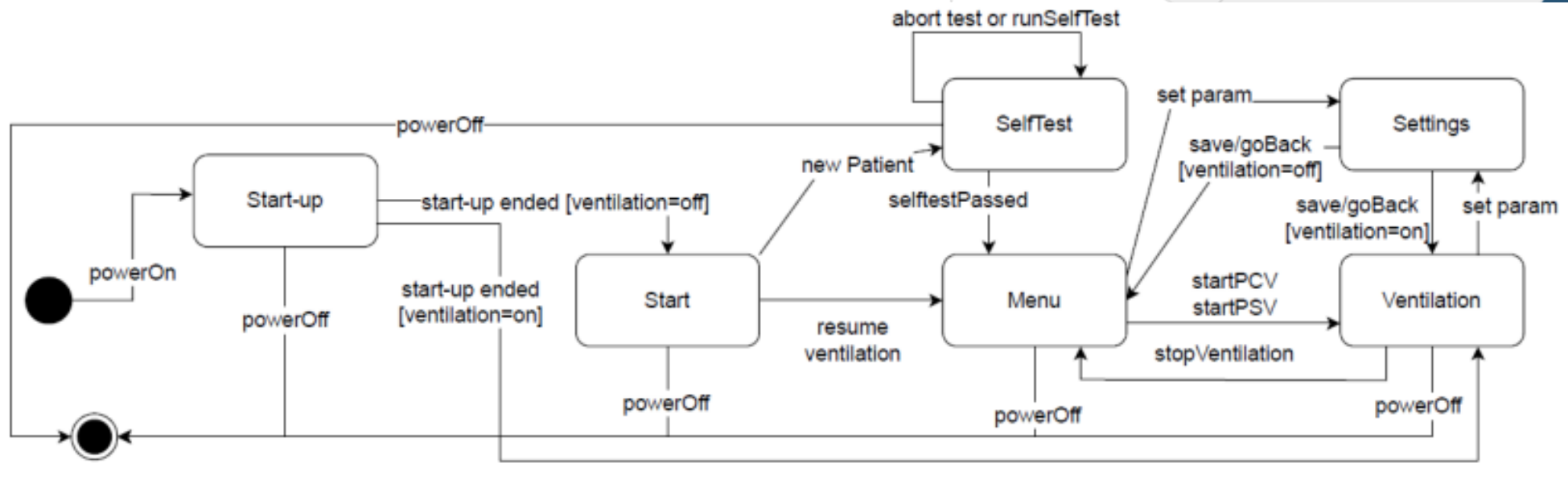


In the contexte C1:

axm1: `partition(Ventilation, {PCV}, {PSV})`

axm2: `partition(ModesG, {StartUp}, {Start}, {Menu}, {SelfTest},
{Settings}, Ventilation, {Off})`

1ST LEVEL: GUI FUNCTIONALITIES (M1+C1)



In the machine M1:

$modeGP \neq modeG \Rightarrow$

$modeGP \mapsto modeG \in$

$possTransG(\mathbf{bool}(power = \mathbf{TRUE} \wedge crashed = \mathbf{FALSE} \wedge$
 $(onAC = \mathbf{TRUE} \vee (switchover = \mathbf{TRUE} \wedge batLev > 0 \wedge batFail = \mathbf{FALSE}))))$

1ST LEVEL: GUI FUNCTIONALITIES (M1+C1)

In the machine M1:

- Transitions with a same label give an Event-B event

Event saveBackAbort $\hat{=}$

any

modeg

where

grd1: $modeG = \mathbf{Settings} \wedge modeg \in ModesG$

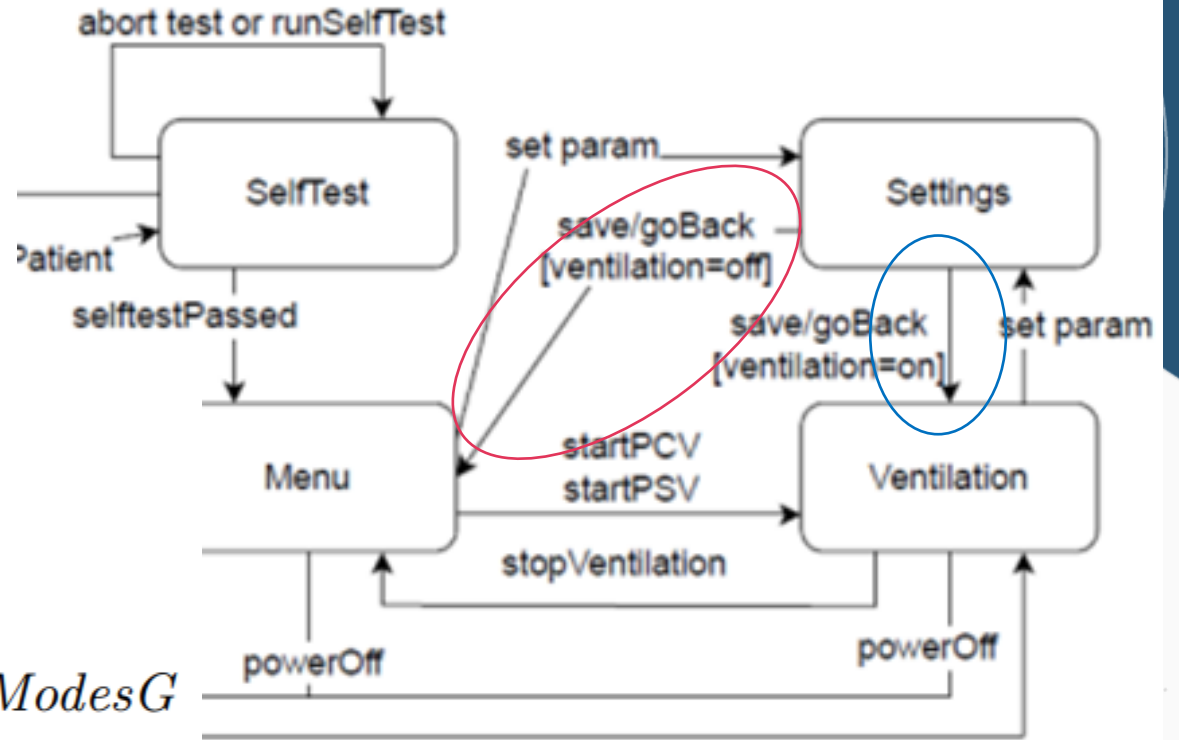
grd2: $modeg \in \mathbf{Ventilation} \cup \mathbf{\{Menu\}}$

then

act1: $modeG := modeg$

act2: $modeGP := modeG$

end



modeGP: previous state of the GUI

modeG: current state of the GUI

modeg: next state of the GUI

2ND LEVEL: CONTROLLER FUNCTIONALITIES

(M2+C2)

In the machine M2:

Event $\text{saveBackAbort} \hat{=}$

any

modeg

where

$\text{grd1: modeG} = \text{Settings} \wedge \text{modeg} \in \text{ModesG}$

$\text{grd2: modeg} \in \text{Ventilation} \cup \{\text{Menu}\}$

$\text{modeC} \neq \text{FailSafe}$

$\text{modeC} \in \text{Ventilation} \Rightarrow \text{modec} \in \text{Ventilation} \wedge \text{modeg} = \text{modec}$

then

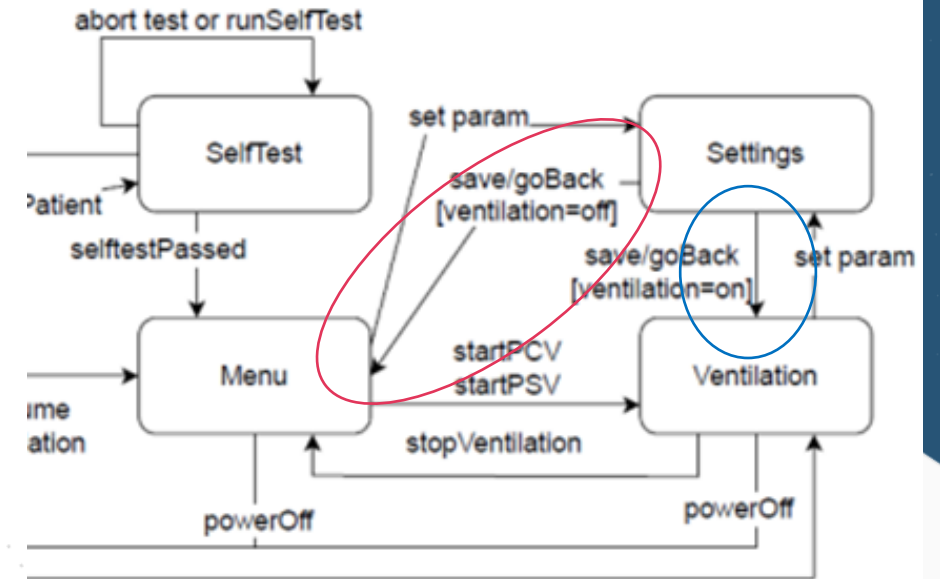
$\text{modeC} \notin \text{Ventilation} \Rightarrow \text{modec} = \text{modeC} \wedge \text{modeg} = \text{Menu}$

$\text{act1: modeG} := \text{modeg}$

$\text{act2: modeGP} := \text{modeG}$

...

end



modeC: current state of the controller
modec: next state of the controller

3RD LEVEL: VENTILATION MODES (M3+C3)

$$psvParamsValC \in 0..curTime \rightarrow (psvParams \rightarrow \mathbb{N}_1)$$
$$modeC = PSV$$
$$\Rightarrow$$
$$\text{dom}(psvParamsValC(\max(\text{dom}(psvParamsValC)))) := psvParams$$

□ **Mode change request**

$$PCV2PSV = \mathbf{TRUE}$$
$$\Rightarrow$$
$$modeG = \mathbf{Settings} \wedge modeC \in \{PCV, \mathbf{FailSafe}\}$$

Event `saveBackAbort` $\hat{=}$ `saveBackAbort`

any

modeg, modec, sv, psvC, psvG, ...

where

`grd1:` ...

...

`grd6:` $PCV2PSV = \mathbf{TRUE} \wedge modeC = PCV \Rightarrow modec = PSV$

`grd7:` $PCV2PSV = \mathbf{FALSE} \vee modeC \neq PCV \Rightarrow modec = modeC$

`grd8:` $psvG = \{\mathbf{TRUE} \mapsto psvParamsValG(\mathbf{max}(\mathbf{dom}(psvParamsValG))),$
 $\mathbf{FALSE} \mapsto psvParamsValC(\mathbf{max}(\mathbf{dom}(psvParamsValC)))\}(sv)$

`grd9:` $psvC = \{\mathbf{TRUE} \mapsto psvParamsValG(\mathbf{max}(\mathbf{dom}(psvParamsValG))),$
 $\mathbf{FALSE} \mapsto psvParamsValC(\mathbf{max}(\mathbf{dom}(psvParamsValC)))\}(sv)$

`grd10:` $modec = PSV \Rightarrow \mathbf{dom}(psvC) = psvParams$

then

`act1:` ...

...

`act5:` $psvParamsValG(curTime) := psvG$

`act6:` $psvParamsValC(curTime) := psvC$

`act7:` $PCV2PSV := \mathbf{FALSE}$

end

4TH LEVEL: VENTILATION PHASES (M4+C4)

In C4: **axm1: partition**(*ventSates*, {*inspBeg*}, {*inspEnd*}, {*expBeg*}, {*expEnd*}, ...)

In M4: **inv1:** *cycleMode* \in *cycles* \rightarrow *Ventilation*

inv2: *ventilPhase* \in *cycles* \rightarrow *ventSates*

inv3: *inspBegT* \in *cycles* \rightarrow \mathbb{N} **inv4:** *inspEndT* \in *cycles* \rightarrow \mathbb{N}

Event *inspStart* $\hat{=}$

any

cy, inspT

where

grd1: *modeC* \in *Ventilation* *cy* \in *Cycles* \setminus *cycles*

...

then

act1: *cycles* := *cycles* \cup {*cy*}

act2: *ventilPhase*(*cy*) := *inspBeg*

act4: *inspBegT*(*cy*) := *curTime*

act5: *inspEndT*(*cy*) := *inspT*

end

Each breathing phase



2 events: *start* and *end*

5TH LEVEL: VALVES (M5+C5)

inv1: $(\exists c. c \in \text{cycles} \wedge \text{ventilPhase}(c) \in \{\text{inspBeg}, \text{inspEnd}\}) \wedge \text{inValveF} = \mathbf{FALSE} \wedge \text{modeC} = \text{Ventilation} \Rightarrow \text{inValve} = \mathbf{TRUE}$

Event $\text{inspStart} \hat{=}$

any

cy, inspT

where

grd1: $\text{modeC} \in \text{Ventilation} \quad cy \in \text{Cycles} \setminus \text{cycles}$

...

then

act1: $\text{cycles} := \text{cycles} \cup \{cy\}$

act2: $\text{ventilPhase}(cy) := \text{inspBeg}$

act4: $\text{inspBegT}(cy) := \text{curTime}$

act5: $\text{inspEndT}(cy) := \text{inspT}$

$\text{inValve} := \{\mathbf{FALSE} \mapsto \mathbf{TRUE}, \mathbf{TRUE} \mapsto \text{inValve}\}(\text{inValveF})$

end

6TH LEVEL: ALARMS (M6+C6)

In C6: `axm1: partition(Alarms, {guiFailure}, {contFailure},
{inValveFailure}, {patConnected}, ...)`

In M6: `alarmRaised ∈ Alarms → BOOL`

6TH LEVEL: ALARMS (M6+C6)

axm1: `partition(Alarms, {guiFailure}, {contFailure},
{inValveFailure}, {patConnected}, ...)`

alarmRaised(inValveFailure) = **bool**($\exists c. (c \in \text{cycles} \wedge$
 $(((((\text{ventilPhase}(c) = \text{inspBeg} \wedge \text{curTime} > \text{inspBegT}(c)) \vee$
 $(\text{ventilPhase}(c) = \text{rmBeg} \wedge \text{curTime} > \text{rmBegT}(c))) \wedge$
 $\text{inValveP} = \mathbf{FALSE}$
 $))) \vee$
 $((\text{ventilPhase}(c) = \text{expPauseBeg} \wedge \text{curTime} > \text{expPauseBegT}(c)) \vee$
 $(\text{ventilPhase}(c) = \text{expBeg} \wedge \text{curTime} > \text{expBegT}(c)) \vee$
 $(\text{ventilPhase}(c) = \text{inspPauseBeg} \wedge \text{curTime} > \text{inspPauseBegT}(c)) \wedge$
 $\text{inValveP} = \mathbf{TRUE}))$))

6TH LEVEL: ALARMS (M6+C6)

Event progress $\hat{=}$
refines progress

any

step, ...

where

grd1: step $\in \mathbb{N}1 \wedge \dots$

grd2: modec $\in \{\text{FailSafe}, \text{modeC}, \text{Off}, \text{StartUp}\}$
...

then

act1: curTime $:= \text{curTime} + \text{step}$

act4: modeC $:= \text{modec}$

end *...*

Event progress $\hat{=}$
refines progress

any

step, ...

where

grd1: $step \in \mathbb{N}1 \wedge \dots$

grd2: $modec \in \{\text{FailSafe}, modeC, \text{Off}, \text{StartUp}\}$

...

$alarmInV = \mathbf{bool}(\exists c. (c \in cycles \wedge$
 $(ventilPhase(c) \in \{inspBeg, inspEnd, rmBeg, rmEnd\} \wedge inValve = \mathbf{FALSE})$
 \vee
 $(ventilPhase(c) \in \{expPauseBeg, inspPauseBeg, inspPauseEnd\} \wedge$
 $inValve = \mathbf{TRUE}))))$

$alarmInV = \mathbf{TRUE} \Rightarrow modec = \text{FailSafe}$

then

act1: $curTime := curTime + step$

act4: $modeC := modec$

$alarmRaised := alarmRaised \triangleleft \{\dots, inValveFailure \mapsto alarmInV\}$

$inValveP := inValve$

...

end

Event progress $\hat{=}$
refines progress

any

step, ...

where

grd1: *step* $\in \mathbb{N}1 \wedge \dots$

grd2: *modec* $\in \{\text{FailSafe}, \text{modeC}, \text{Off}, \text{StartUp}\}$

...

$\text{alarmInV} = \mathbf{bool}(\exists c. (c \in \text{cycles} \wedge$
 $(\text{ventilPhase}(c) \in \{\text{inspBeg}, \text{inspEnd}, \text{rmBeg}, \text{rmEnd}\} \wedge \text{inValve} = \mathbf{FALSE})$
 \vee
 $(\text{ventilPhase}(c) \in \{\text{expPauseBeg}, \text{inspPauseBeg}, \text{inspPauseEnd}\} \wedge$
 $\text{inValve} = \mathbf{TRUE}))))$

$\text{alarmInV} = \mathbf{TRUE} \Rightarrow \text{modec} = \mathbf{FailSafe}$

then

act1: *curTime* := *curTime* + *step*

act4: *modeC* := *modec*

alarmRaised := *alarmRaised* $\Leftarrow \{\dots, \text{inValveFailure} \mapsto \text{alarmInV}\}$

inValveP := *inValve*

...

end

WHAT HAS BEEN DONE ...

- Different states/transitions of the controller and the GUI
- Ventilation modes (PCV and PSV) and the switching from one to the other
- Ventilation parameters and their update before and during the ventilation
- Position of the valves (in and out) during the ventilation and their failures
- Alarms: the valves (in and out), patient connection while the system is in the state StartUp, the backup battery, the switchover, the FI1/FI12/oxygen sensors.

SOME STATISTICS

- ❑ Three months development/proof
 - 6 refinement levels
 - 44 variables + 90 invariants +35 events
- ❑ Models are proved and verified
 - Model checking: ProB for detecting obvious invariant violations
 - Validation: ProB on our own scenarios
 - Proof obligations correctness
 - 1322 proof obligations
 - 312 automatic (23%)
 - 225 interactive

FEEDBACKS ON THE SPECIFICATION DOCUMENT

- Under which conditions the controller must move to the *state FailSafe*?
- Does the controller continuously check the presence of undesirable events or not?
- Does the controller continuously check the communication with the GUI or only in the state *StatUp*?
- What is the impact of the alarms' levels on the system?